Programming Paradigms CT331 Week 8 Lecture 2

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Binary Search Tree (BST)

Data structure to store an ordered set of items.

- Fast lookup
- Fast addition and removal of items
- Can be used to implement:
 - Dynamic sets of items
 - Lookup tables that allow finding an item by its key



Binary Search Tree rules

Three rules for all BSTs:

- 1. No duplicates
- 2. Maximum of 2 child nodes.
- Left child is always smaller than parent. (and right is always larger)

Each child, or leaf node is also a BST.

We call these sub-trees.

Smallest possible sub-tree: null

Binary search trees keep their keys in sorted order, so that lookup and other operations can use the principle of binary search.

When looking for a key in a tree (or a place to insert a new key):

Traverse the tree from root to leaf, making comparisons to keys stored in the nodes of the tree and deciding, based on the comparison, to continue searching in the left or right subtrees.

On average, this means that each comparison allows the operations to skip about **half of the tree**.

- Much better than linear search through unsorted list
- Not as fast as a hash table operations.

Binary Search Trees: Operations

- Searching
- Insertion
- Deletion
- Traversal
- Verification

Use the following slides as algorithms for the assignment

Find the node with value 6:

Start at head (8)

- 1. Is current node = null?
 - a. Y: item not in tree
 - b. N: continue
- 2. Is current node = 6?
 - a. Y: Return, item found.
 - b. N: Continue
- 3. Is 6 < current node?
 - a. Y: Recurse on left child/leaf
 - b. N: Recurse on right child/leaf



Insert the value: 5

Start at head (8)

- 1. Is current node = null?
 - a. Y: set current node to 5, return.
 - b. N: continue
- 2. Is current node = 5?
 - a. Y: Return, item found.
 - b. N: Continue
- 3. Is 5 < current node?
 - a. Y: Recurse on left child/leaf
 - b. N: Recurse on right child/leaf



Delete the node with value 3:

Start at head (8)

- 1. Is current node = null?
 - a. Y: item not in tree
 - b. N: continue
- 2. Is current node = 3?
 - a. Y: Return, item found.
 - b. N: Continue
- 3. Is 3 < current node?
 - a. Y: Recurse on left child/leaf
 - b. N: Recurse on right child/leaf



Deletion is not so simple...

- If node has no children:
 - Just set to null.
- If node has 1 child:
 - Set node value to child value
 - Delete child.
- If node has 2 children:
 - o **???**



If node has 2 children:

1. Find in-order successor (IOS).

(the next highest value in the tree) (the smallest value on the right sub-tree)

2. Copy IOS to node

Delete old IOS.

Note: you can also use the in-order predecessor, or the next lowest value, the highest value on the left sub-tree.



Once the binary search tree has been created, its elements can be retrieved in-order by:

- Recursively traversing the left subtree of the root node.
- Accessing the node itself
- Then recursively traversing the right subtree of the node



How do we represent a BST in scheme?

Minimum sub-tree:

Null or '()

Typical sub-tree:

(Left child, value, right child)



How do we get the right child?

How do we get the value?

(car sub-tree)

How do we get the right child?



How do we get the value?

(car sub-tree)

How do we get the right child?

(caddr sub-tree)

How do we get the value?



(car sub-tree)

How do we get the right child?

(caddr sub-tree)

How do we get the value?

(cadr sub-tree)

