#### CT420 REAL-TIME SYSTEMS

# DESIGN CONSIDERATIONS REAL-TIME SAFETY-CRITICAL SYSTEMS

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# Motivation for and Objectives of this Lecture

- So far, we have addressed the real-time aspects of RTSCS
  - How must software systems be designed to meet RT requirements?
  - What APIs are available (i.e., POSIX)?
- In this lecture we focus on features that increase system safety by increasing its reliability

#### Recall Case Study: The Boeing 737 Max 8 MCAS Problem

- On 29 October 2018, a just 2-months old Boeing 737 MAX 8 plane crashed into the Java Sea 12 minutes after takeoff, killing all 189 passengers and crew
- The plane's flight recorder showed the following vertical speed pattern before the crash:



- Then on March 10, 2019, a 4-months old 737 MAX 8 crashed shortly after take-off from Addis Ababa, killing all 149 passengers and 8 crew members on board
  - Evidence retrieved later suggested that again the aircraft's vertical speed after take-off was unstable
- □ Shortly later, the entire fleet was grounded

#### The Boeing 737 Max 8

- The latest and most fuel-efficient version of the Boeing 737
  - The 737 series is the highest-selling commercial jetliner in history with more than 10,000 units built since 1967
  - The Max 8 has longer engines than previous models, which sit slightly forward and higher, therefore changing its centre of gravity, making it more likely to pitch upward on take-off





#### Boeing's Manoeuvring Characteristics Augmentation System (MCAS)



Source: <u>https://theaircurrent.com/aviation-safety/what-is-the-boeing-737-max-maneuvering-characteristics-augmentation-system-mcas-jt610/</u>

#### Boeing's Manoeuvring Characteristics Augmentation System (MCAS)



an approaching stall ...

flight

Reporting by DOMINIC GATES, Graphic by MARK NOWLIN / THE SEATTLE TIMES



Source: <u>https://www.seattletimes.com/business/boeing-aerospace/failed-certification-faa-missed-safety-issues-in-the-737-max-system-implicated-in-the-lion-air-crash/</u>

### The Angle of Attack (AoA) Sensor

- In both crashes the incorrect airflow angles, reported from only a single (faulty) AoA sensor, were processed
- Using multiple AoA sensors would have allowed to compensate for this
  - Note that the plane had in fact 2 AoA sensors installed
- □ Compare to Airbag design we discussed before





#### Other Issues that led to the Crashes

 MCAS should not be able to repeatedly overwrite pilot decisions



- Pilots should have been trained how to manually disable the MCAS in-flight <u>http://www.spiegel.de/wissenschaft/technik/boeing-737-max-abstuerze-welche-rolle-spielten-die-piloten-a-1258835.html</u> (German article)
- 3. The Lion Air machine did not have dashboard instruments to show both AoA readings, or to alert the pilots about a discrepancy These were optional extras Lion Air did not want to pay for http://www.spiegel.de/wissenschaft/technik/boeing-737-maxfehlten-sicherheitsfunktionen-weil-sie-extra-kosten-a-1259117.html (German article)

#### Recap: Quality Requirements for RTSCS

- RTSCS must be time responsive
- RTSCS must be reliable
  - The ability to behave in accordance with its specification
- RTSCS must be safe

Conditions that lead to hazards do not occur

- RTSCS must be secure
  - Protect itself against intentional or accidental access, use, modification or destruction
- RTSCS must be usable
  - Easy to learn, understand, and use
- RTSCS must be maintainable
  - Return swiftly to an operational state after receiving repairs or modification (e.g. plug in-and-forget)

#### Accident, Risk and Hazard

- Accident is a loss of some kind, such as injury, death, or equipment damage
- Risk is a combination of the likelihood of an accident p(a) and its severity s(a):
  - Often numerical models are used with p(a) being based on a probability distribution (i.e. 0 <= p(a) <= 1), and s being or normed value (i.e. 0 <= s(a) <= 1): risk = p(a) \* s(a)
- Hazard is a set of conditions and/or events that leads to an accident

#### Faults and Hazards

- Faults lead to hazards, which lead to accidents
- Faults are the manifestation of a:
  - Failure
    - Random non-performance of a component (e.g. wear and tear)
  - Error
    - Systematic; i.e. design fault or software fault (bug)
- Faults can be permanent, intermittent, or transient

#### **Determining Failure Probability**



- Assume a driver airbag system with N (= 5) independent components, each with a fault probability of 5% over 10 years of operation (i.e. 95% probability that component will function ok after 10 years)
- Overall system failure probability after 10 years (assuming statistical independence):
  - 1 (Probability that ALL components ok)
  - **1** 1-  $(1 0.05)^5$  = 1 0.773 = 0.227 = ~22.7%
- Components may include
  - Sensors, actuators, controllers
  - Their components, e.g. CPU, RAM, storage

#### The Bathtub Curve

- The bathtub curve is a particular shape of a failure rate graph
  - Failure rate is the frequency with which an engineered system or component fails over time
- Component failures can be pre-empted via
  - Component redundancy
  - Scheduled component replacement



#### Fault Tree Analysis (FTA)

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- FTA is a top-down (from hazard / event to basic fault type), deductive failure analysis in which an undesired state of a system is analysed using Boolean logic to combine a series of lower-level events
- This analysis method is mainly used in safety engineering and reliability engineering to understand how systems can fail, to identify the best ways to reduce risk and to determine (or get a feeling for) event rates of a safety accident or a particular system level (functional) failure

## Fault Tree Analysis Symbology

An event that results from a combination of events through a logic gate

A basic fault event that requires no further development

A fault event because the event is inconsequential or the necesary information is not available

An event that is expected to occur normally



A condition that must be present to produce the output of a gate



Transfer



AND gate



**OR** Gate



NOT Gate

#### **Example Pacemaker**

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### Example Pacemaker (Subset) Fault Analysis



### Making RTSCS safe: Fail-Safe

- Fail-safe describes a device or feature which, in the event of failure, responds in a way that will cause no harm or at least a minimum of harm to other devices or danger to personnel
- Examples:
  - Air brakes on railway trains and air brakes on trucks
  - Luggage carts in airports
  - Lawnmowers

### Making RTSCS safe: Fail-Soft

- Pertaining to or noting facilities built into a system, as in an automobile or a computer, for continuing operations on an <u>interim basis</u> and probably with reduced efficiency, if parts of the system fail
- Example: Fail-Soft of ECU via "Limp Mode"





#### Making RTSCS safe: Graceful Degradation

- As size of faulty set increases, system **must** not suddenly collapse, but must gracefully degrade
- □ Failures will eventually impact on (RTS) operation
- System perhaps operates with reduced functionality
- Avoid catastrophic failure



# Example of graceful Degradation: The Citroen CX

- Common hydraulic
   system for steering,
   brakes and suspension
- What goes first, second and last when hydraulic pressure drops?



#### Fault Types

#### Permanent

- Easiest to detect
- Hardware failure or software design/code fault

#### Intermittent

- Fault appears from time to time
- Loose wire, poor contacts, certain sequence of events

Transient

- Fault appears but dies away with time
- **α** particle impact: non-destructive to memory

#### Transient Soft Errors due to Neutron Strikes or α Particle Impacts





- The capacitor of a DRAM cells stores a single bit (0 = no charge, 1 = charged), while high DRAM integration density result in very small capacitors
- An α particle penetrating a DRAM cell may have sufficient kinetic energy to distort the capacitor charge and changes its state (chip-level soft error)
- α particles can be emitted from radioactively contaminated packaging
- □ As a result, high-end servers use DRAM with error detection/correction capabilities (→ information redundancy)

#### Information Redundancy

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- Data Validity Checks of data at rest (i.e., in RAM, or in secondary storage) can be achieved via:
  - Cyclic Redundancy Check (CRC)
    - Here blocks of data get a 16-bit or 32-bit check value attached, based on the remainder of a polynomial division of their contents
    - This identifies
      - all single or dual bit errors
      - a high percentage of multiple bit errors
    - Widely used in data communication (e.g, TCP/IP)
  - One's complement
  - Error correcting codes
  - RAID
- Redundancy should be set every write access
- Data should be checked every read access

#### **One's Complement Data Validity Check**



### Error Correcting Codes (ECC)

- ECC is redundant information added to the data / message
- This allows detecting and correcting a limited number of errors that may occur anywhere in the message
- The American mathematician Richard Hamming pioneered this field in the 1940s and invented the first error-correcting code in 1950: the Hamming (7,4) code

#### Hamming Distance

- Minimum number of bits to be toggled to convert one codeword into another
  - □ HD (ASCII) = 1
  - HD (ASCII + Parity Bit) = 2
    - even parity: Total number of "1"s is <u>even</u> number; odd parity: ditto
- $\Box$  A code with a given HD x may be able to
  - detect (x 1) bit errors
  - correct (1 <= y < x / 2) bit errors
- Examples:
  - ASCII
    - "@"= $64_{10} = 01000000$
    - "A"=65<sub>10</sub> = 01000001
    - Hamming Distance 1
  - ASCII + even parity
    - "@"=64<sub>10</sub> = 01000000 1
    - "A"= $65_{10}$  = 01000001 0
    - Hamming Distance 2

## Hamming (7,4) Code

Hamming (7,4) is an error-correcting code that encodes four bits of data into seven bits by adding three (even or odd) parity bits

| Bit #           | 1     | 2     | 3     | 4     | 5     | 6     | 7     |
|-----------------|-------|-------|-------|-------|-------|-------|-------|
| Transmitted bit | $p_1$ | $p_2$ | $d_1$ | $p_3$ | $d_2$ | $d_3$ | $d_4$ |
| $p_1$           | Yes   | No    | Yes   | No    | Yes   | No    | Yes   |
| $p_2$           | No    | Yes   | Yes   | No    | No    | Yes   | Yes   |
| $p_3$           | No    | No    | No    | Yes   | Yes   | Yes   | Yes   |



# Example: Hamming (7, 4) Code with even Parity



| d1  | d2 | d3 | d4 | pl | p2       | р3       | Start                                  |
|-----|----|----|----|----|----------|----------|--|
| 1   | 1  | 0  | 1  | 1  | 0        | 0        | Juli                                   |
|     |    |    |    |    |          |          |  |
| d 1 | d2 | d3 | d4 | pl | p2       | р3       | Correction of single<br>data bit error |
| 0   | 1  | 0  | 1  | 1  | <u>0</u> | 0        |  |
| d 1 | d2 | d3 | d4 | pl | p2       | р3       | Correction of single                   |
| 1   | 1  | 0  | 1  | 1  | <u>0</u> | 1        | parity bit error                       |
| d 1 | d2 | d3 | d4 | pl | p2       | р3       | Detection of double                    |
| 0   | 1  | 0  | 0  | 1  | 0        | <u>0</u> | data bit error                         |

Hamming(7,4) has a HD of 3!

#### Problem

- □ Consider a Hamming (7, 4) code with even parity
- Does the Bitvector "1101: 010" (d1 d4 : p1 p3) indicate a bit-error? If yes, which bit got flipped?
- Draw a Venn Diagram to figure it out...



#### Hamming(7, 4) Problem Solution



#### Mass-Storage Redundancy via RAID

Redundant Array of Independent Disks (RAID) is a data storage virtualisation technology that combines multiple physical disk drive components into one or more logical units for the purposes of data redundancy and / or performance improvement



- Data blocks are distributed across the drives in one of several across the drives in one of severa
- Many RAID levels use a parity-based error protection scheme (see RAID-4), example (with 12 bit / block):
  - Block 1: 010001101001
  - Block 2: 110011011010
  - Block 3: 000100100101
  - P Block: 100111010110 (bitwise EXOR, equivalent to even parity)

#### Mass-Storage Redundancy via RAID

- RAID storage systems require a dedicated RAID controller, that supports the required RAID level
   See also the diagram on the next slide
  - Normally such controllers are not shown in RAID diagrams



- Block-level striping without parity or mirroring
  - data striping is the technique of segmenting logically sequential data, such as a file, so that consecutive segments are stored on different physical storage devices
- 2 or more drives (n) required
- No redundancy, but up to n-times R/W performance increase



- Block-level mirroring without parity or striping
- 2 or more drives (n) required
- (n 1) drive failures can be compensated; here each disk can
  - diagnose catastrophic failures (e.g. head crash)
  - detect (but not correct) sector-wise bit errors on platters
- □ No increase in R/W performance





- Block-level striping with single parity disk
- Single catastrophic drive failure can be compensated (<u>any</u> drive can fail)
- RAID 4 provides good performance of random reads, while the performance of random writes is low due to the need to write all parity data to a single disk (Disk 3 in the diagram above)
- □ Minimum of 3 drives required



#### Drive Hot-Swapping in RAID

#### □ In RAID a defect drive will be (ASAP)

- manually swapped for a new drive (hot-swap), or
- replaced by an idle drive (hot-spare) already in the system
- The new drive's content is rebuild by the RAID controller while the disk set is still operational
  RAID 4
- Example RAID 4 with Disk 0 swapped:
  - A1 = A2 EXOR A3 EXOR  $A_P$
  - **B**1 = B2 EXOR B3 EXOR  $B_P$
  - $C1 = C2 EXOR C3 EXOR C_P$
  - **D** $1 = D2 EXOR D3 EXOR D_P$



- □ Similar to RAID 4, but:
  - Block-level striping with distributed parity
  - Distributed parity evens out the stress of a dedicated parity disk among all RAID members
  - Write performance is increased since all RAID members participate in the serving of write requests
- □ Minimum of 3 drives required



#### Increasing Hardware Reliability

- Information redundancy protects against memoryrelated faults
- However, it assumes that the underlying computer system works satisfactorily
- □ Therefore, we need to determine ways for a system's
  - Fault detection
    - $\rightarrow$  Watchdog
  - Fault recovery
    - $\blacksquare 
      ightarrow$  Failover
    - $\blacksquare \rightarrow$  Redundancy

#### Watchdog

- Idea: Restart can be a fail-safe state!
- Here the computer / CPU is reset by the watchdog, unless it is regularly (typically every 1ms to every 10s) triggered by the CPU, for example using a PIO
- Requires validation that the system is in safe state during reset
- Watchdog can be internal or external (as shown) component





#### Example for internal Watchdog (Arduino Uno)





#define wdt\_reset(); //resets WDT
#define wdt\_disable(); //disables WDT
#define wdt\_enable(timeout); //sets the
watchdog pre-scaler, using one of the
constants below:

#define WDTO\_15MS 0
#define WDTO\_30MS 1
#define WDTO\_60MS 2
#define WDTO\_120MS 3
#define WDTO\_250MS 4
#define WDTO\_500MS 5
#define WDTO\_1S 6
#define WDTO\_2S 7
#define WDTO\_4S 8
#define WDTO 8S 9

#### Example for internal Watchdog (Arduino Uno)

#include <avr/wdt.h>
void setup(){
 Serial.begin(9600);
 Serial.println("Setup started :");
 delay(2000);
 wdt\_enable(WDTO\_4S);
}
void loop(){

```
Serial.println("LOOP started ! ");
for(int i=0; i<=5; i++){
    Serial.print("Loop : ");
    Serial.print(i);
    Serial.println();
    delay(1000); }
    wdt_reset();</pre>
```

```
//infinity loop to hang MCU
while(1){}
```

}

| COMM                      | (2 <del>55)</del> -             |              | M     |
|---------------------------|---------------------------------|--------------|-------|
|                           |                                 | Send         | · · · |
| Setup started :           |                                 |              |       |
| GOOP started !            |                                 |              |       |
| 600p : 0                  |                                 |              |       |
| loop : 1                  |                                 |              |       |
| loop : 2                  |                                 |              |       |
| loop : 3                  |                                 |              |       |
| 600p : 4                  |                                 |              |       |
| loop : 5                  |                                 |              |       |
| Setup started :           |                                 |              |       |
| GOOP started !            |                                 |              |       |
| loop : 0                  |                                 |              |       |
| Loop : 1                  |                                 |              |       |
| 100p : 2                  |                                 |              |       |
| loop:3                    |                                 |              |       |
| 600p : 4                  |                                 |              |       |
| 600p : 5                  |                                 |              |       |
| Setup started :           |                                 |              |       |
| LOOP started !            |                                 |              |       |
| loop : 0                  |                                 |              |       |
| loop : 1                  |                                 |              |       |
|                           |                                 |              |       |
|                           |                                 |              |       |
| Autoscroll Show timestamp | No line ending \vee 9600 baud 🗸 | Clear output |       |

### Watchdog Coding Challenge

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- Consider the Arduino controls a chemical reactor in a factory, where it is continuously operating. A watchdog monitors the Arduino
  - The control code is executed in loop(){}
- Every 3 months the reactor is turned off and serviced
- Here technicians need to establish, how many watchdog resets occurred since the last service
  - A watchdog reset occurs if the code within a single loop iteration is not completed in time
- □ This value is stored in some non-volatile (Flash) memory
- This memory can store 32 integer values, and can be accessed via
  - int flash\_read(int address) // address is between 0 and 31
  - flash\_write(int address, int value) // ditto
- The Flash memory is reset (all Os) when the system is deployed first and after each service
- Use as much as needed for your solution
- Hint: Use both the setup() and loop() routine



#### My Solution

...

```
loop() {
 flash_write(0, 1); // write a "1" to Flash cell #0
  // Execute loop code
  ...
  flash_write(0, 0); // write a "0" to Flash cell #0
 wdt_reset();
}
setup() {
 if (flash_read(0) == 1) \{// WDT reset occurred
    flash_write(1, flash_read(1) + 1); // increment counter
    flash_write(0, 0) // Reset flag
```

| Flash Memory Cell #0    | Flash Memory Cell #1 |
|-------------------------|----------------------|
| 1 == WDT reset; 0 == Ok | #of WDT resets       |

#### Master-Slave Fail Over



- A watchdog reset may disable a single controller for too long, therefore a second controller in stand-by mode may take over instead
- Here the equivalent of a watchdog reset pulse (the Alive signal) send by the active computer A is monitored by the passive computer B
- □ If computer A fails to provide this pulse in time, a timeout will occur, causing computer B to take over, while keeping computer A in a reset state

#### Redundancy via Synchronously Operating and Clocked independent Computers

- Example Moneypoint's Burner
   Management System Siemens
   AS220
  - Complex control system that manages power plant
- Triple-redundant hardware
  - RAM, ROM, CPU
  - CPUs run in synchronously
- 2-out-of-3 voters are used to deal with single faulty component



#### AS220 EHF Mode Of Operation



#### Increasing Sensor Reliability

- The MCAS example showed that malfunctioning system sensors can lead to hazards and accidents
- □ Therefore, multiple redundant sensors must be put in place
- The (MCAS) computer detects the sensor reading discrepancy
- □ But...
  - What degree of difference indicates a faulty unit?
    - This has very much to be decided on a case-by-case basis
  - How can one identify the faulty unit in the first place?

#### NMR: N-Modular Redundancy (Single Voter)



- Simple Design based on identical sensors P1, P2 and P3
- The voter determines if a sensor is faulty, i.e. returns incorrect readings
- Good for dealing with random faults only
- Only feasible if voter has a much lower failure probability than P1, P2 and P3

# Example NMR (Single Voter)

- N identical independent components, each failure probability of 5%
- □ 2v3 System → 2 units need to fail for a complete system failure
- Overall system reliability: 1- (0.05)<sup>2</sup>
   = 99.75%
- Compared to 95% system reliability without redundancy
- This of course assumes that the voter failure probability is negatable small
- How can such a voter be implemented?



## Voter Types

- All voter types try to determine a correct sensor reading
- The top 3 voters may be used to identify faulty units
- Formalised majority voter (FMV)
  - All inputs are equal, selects absolute (> 50%) majority
- Generalised median voter (GMV)
  - All inputs are equal, selects the median of the values
- Formalised plurality voter (FPV)
  - All inputs are equal, partitions the set of inputs based on metric equality and selects the output from the largest group, i.e. picks most common value
- Weighted averaging (WA)
  - Combines the outputs in a weighted average (mean)

1, 3, 3, **6**, 7, 8, 9 Median = <u>6</u> 1, 2, 3, **4**, **5**, 6, 8, 9 Median = (4 + 5) ÷ 2 = <u>4.5</u>

#### Case Study

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- Consider a rocket engine as shown in the diagram
- Its high-pressure fuel turbopump has to operate within a certain fuel pressure range that is constantly monitored by a range of pressure sensors (not shown)
- Because of the extreme environment the sensor readings have a significant error, so the readings or multiple sensors are processed using a voter
- The diagram on the next slide shows a set of readings (in Megapascal) provided by 5 pressure sensors
- The sensors have a different weight depending on their perceived reliability



# Calculate FMV, GMV, FPV and WA Output for the Data below

| Tag    | Input 1 | Input 2 | Input 3 | Input 4 | Input 5 |
|--------|---------|---------|---------|---------|---------|
| Weight | 0.2     | 0.3     | 0.1     | 0.3     | 0.1     |
| Val    | 10      | 11      | 12      | 11      | 8       |
| FMV    |         |         |         |         |         |
| GMV    |         |         |         |         |         |
| FPV    |         |         |         |         |         |
| WA     |         |         |         |         |         |

- Calculate FMV, GMV, FPV and WA output for the data above
- Do not round results, i.e. return decimal values where appropriate
- Note that the weight row adds up to 1.0

# Calculate FMV, GMV, FPV and WA output for the Data below

| Tag    | Input 1           | Input 2 | Input 3 | Input 4 | Input 5 |  |  |
|--------|-------------------|---------|---------|---------|---------|--|--|
| Weight | 0.2               | 0.3     | 0.1     | 0.3     | 0.1     |  |  |
| Val    | 10                | 11      | 12      | 11      | 8       |  |  |
| FMV    | n/a               |         |         |         |         |  |  |
| GMV    | 11                |         |         |         |         |  |  |
| FPV    | 11                |         |         |         |         |  |  |
| WA     | 10.6 (rounded 11) |         |         |         |         |  |  |

- Calculate FMV, GMV, FPV and WA output for the data above
- Do not round results, i.e. return decimal values where appropriate
- Note that the weight row adds up to 1.0

#### Increasing Software Reliability

- System faults may not only be the result of hardware issues, but can be a consequence of software problems
- Again, redundancy approaches can help, i.e.
  - Static software redundancy
  - dynamic software redundancy
- Let's start with a case study first (Ariane 5)

#### The Ariane Rocket Family

- Ariane is a series of a European civilian (ESA) expendable launch vehicles for space launch use
- GTO =
   geosynchronous
   transfer orbit



#### The Ariane 5 Accident

- Ariane 5 is a now retired European heavy-lift space launch vehicle
- The launch vehicle had 82 consecutive successful launches between 2003 and 2017
- However, its maiden flight on 4 June 1996 resulted in self-destruction after 37 seconds because of a malfunction in the control software
- See
  <u>https://www.youtube.com/watch?v=gp</u>
  <u>D8r-2hwk</u>



#### One Bug – One Crash

- Steering was controlled by the on-board computer, which mistakenly thought the rocket needed a course change because of numbers coming from the inertial guidance system (IGS). The IGS uses gyroscopes and accelerometers to track motion The numbers looked like bizarre flight data, but were actually a diagnostic
  - error message. The guidance system had in fact shut down!
- This shutdown occurred 36.7 seconds after launch, when the guidance system's own computer tried to convert one piece of data the sideways velocity of the rocket from a 64-bit format to a 16-bit format. The number was too big, and an overflow error resulted. When the guidance system shut down, it passed control to an identical, redundant unit, which was there to provide backup in case of just such a failure. But the second unit had failed in the identical manner a few milliseconds before, as it was running the same software

#### Ariane 5 Flight Controller



#### The Ariane 5 Accident: Root Cause Analysis

- The software for the IGS was originally written for the Ariane 4 rocket and re-used for Ariane 5
- However, Ariane 5 is a more powerful rocket than Ariane 4, resulting in a sideways velocity that was not anticipated when the IGS was originally build, causing a numeric overflow
- As a result, the booster nozzles got incorrectly aligned by the on-board computer, which led to a rapid change of attitude, which caused the launcher to disintegrate due to aerodynamic forces





#### Software Redundancy

The most certain and effectual check upon errors which arise in the process of computation is to cause the same computations to be made by separate and independent computers, and this check is rendered still more decisive if their computations are carried out by different methods

#### Static Software Redundancy

#### □ N version programming (N>=2)

- Independent generation of N functionally equivalent programs from same spec

#### N versions run concurrently

- Voter makes decision
- Impacts on performance and synchronisation issues may have to be considered
- Good match for NMR
- Also
  - Use of a common language may lead to common errors
  - Different compilers/hardware minimise risk of common failure

#### **N-Version Programming**



#### N-Version Programming: Example Factorial



#### N-Version Programming Issues

Redundant code runs regardless of whether faults are present

- Additional CPU-resources required
- Redundant code may have different execution times, may cause synchronisation problems
- Initial spec may provide common failure mode regardless of subsequent strategies
- This requires the creation of diverse and equivalent specifications so that programmers can design software which do not share common faults
- Teams of programmers may have similar bad habits and/or biased programming techniques
  - Especially if same language is used (think of pointers in C)
- □ Overall, very costly process → limited application (e.g. avionics, military)

#### Dynamic Software Redundancy

□ In this approach, redundant components run only under fault conditions

- This assumes that a fault can be detected, e.g. via overflow, out-of-bounds, or acceptability checks
- □ After a fault has been detected, **backward recovery** takes place:
  - Go back and restore system to safe state prior to error
  - Avoid problem on repeat by using different implementation

#### **Recovery Block Approach**



#### Recovery Block Approach via Java Exception Handling

```
Element getElement(x) {
    try {
        return repository.getElement(x);
    } catch (NotFoundException e) {
        return fallbackToSimilar(x);
Element fallbackToSimilar(x) {
    try {
        return repository.getSimilarElement(x);
     } catch (NotFoundException e1) {
        return fallbackToParent(x);
Element fallbackToParent(x) {
    try {
        return repository.getParentElement(x);
    } catch (NotFoundException e2) {
        throw new IllegalArgumentException(e);
```

- Consider three versions, i.e.
  - getElement()
  - fallbackToSimilar()
  - fallbackToParent();
- If getElement() throughs a NotFoundException object, fallbackToSimilar() will be called
- If fallbackToSimilar() throughs a NotFoundException object, fallbackToParent() will be called

#### Summary

- Hardware-, software-, and information redundancy are the building blocks for RTSCS
- They increase system reliability and subsequently system safety
- They can be found in many industries including
  - Automation
  - Avionics (next slides)
  - Military
  - Medical device
  - Robotics

#### Final Example: Airbus 340

#### Hardware & software redundancy

- NMR redundancy with
  - 3 main flight controllers
  - 2 backup flight controllers (that replace faulty unit on-the-fly)
- Software developed by different teams and on different platforms





#### Final Example: Boeing 777

- Hardware Redundancy
  - Motorola
  - AMD
  - Intel
    - Example <u>Pentium FDIV bug</u>
- Ada programming language used, but different compilers



The correct value is:  $\frac{4,195,835}{3,145,727} = 1.333820449136241002$ When converted to the hexadecimal value use at or beyond four digits:<sup>[9][10]</sup>  $\frac{4,195,835}{3,145,727} = 1.333739068902037589$