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## CS4423-Networks: Week 11 (26+27 March 2025)

## Part 1: Watts-Strogatz model

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This Jupyter notebook, and PDF and HTML versions, can be found at https://www.niallmadden.ie/2425-CS4423/#Week11

This notebook was adapted by Niall Madden from one developed by Angela Carnevale.

Our usual preamble:

```
In [1]: import networkx as nx
import numpy as np
opts = { "with_labels": True, "node_color": "#84003d", "font_color": "white" } # Galu
import random # some random number generators:random, random_choices
import statistics # e.g., mean of entries in a list
import math # for comb (=binomial coef)
import matplotlib.pyplot as plt
np.set_printoptions(precision=2) # just display arrays to 2 decimal places
np.set_printoptions(suppress=True)
```

## Small Worlds, again

Last week, we claimed that **small world networks** tend to share three characteristics:

- 1. Short **characteristic path length**, which scales like  $\ln n$ , where n is the number of nodes.
- 2. Low transitivity, meaning that a high proportion of triads form triangles.
- 3. A high clustering coefficient

We saw that the  $G_{ER}$  models tend to have the first property, but not the second or third. Therefore,

in this sense, they don't mimicking real world networks very well.

An alternative, developed by Watts and Strogatz in 1998, is to start with some **regular network** that naturally has a **high clustering**, and then to randomly distort its edges, to introduce some **short paths**.

## **Circle Graphs**

#### Cycle graphs

To learn how to create a network that has the right properties, we'll start with one that does not, and then see how we can change to.

So we start with a **cycle graph**.

```
In [2]: n = 16
G = nx.cycle_graph(n)
nx.draw_circular(G, **opts)
print(f"G has an average path length of L={nx.average_shortest_path_length(G):.3}")
print(f"Its transitivity value is T={nx.transitivity(G)}, and Clustering is C={nx.ave
C has an average path length of L=4.27
```

G has an average path length of L=4.27 Its transitivity value is T=0, and Clustering is C=0.0



We won't dwell on it right now, but the average path length is pprox n/4. Let's focus of the transitivity and clustering.

• **Transitivity:** clearly,  $C_n$  has many triads, but no triangles.

• **Clustering**: the subgraph induced by the neighbours of any node has no edges.

To address this, we could add edges between nodes that have a common neighbour.

#### Increasing Clustering

Starting with the cycle graph,  $G = C_n$ , let's add an edge between Node i and Node i + 2 (mod n).

```
In [3]: for v in G:
    G.add_edge(v, (v+2) % n)
    nx.draw_circular(G, **opts)
    print(f"For this graph, G, we have L={nx.average_shortest_path_length(G):.3}, T = {nx
    For this graph, G, we have L=2.4, T = 0.5, and C = 0.5
```



Looks like we're going in the right direction: L is getting smaller while C (and T) are increasing. Let's keep going by adding an edge between Node i and Node i + 3 (mod n).

In [4]:	for v in G:
	G.add_edge(v, (v+3) % n)
	<pre>nx.draw circular(G, **opts)</pre>
	<pre>print(f"For this graph, G, we have L={nx.average_shortest_path_length(G):.3}, and C =</pre>
F	For this graph, G, we have L=1.8, and C = $0.6$



#### **Circle Graph: Definition**

**Definition (Circle Graph).** For 1 < d < n/2, an (n, d)-**circle graph** is obtained from a cycle on n vertices by additionally linking each node to all nodes that are not more than d steps away on the cycle.

Here is some code to generated it:

```
print(f"For this graph, G, we have L={CPL:.3f}, C={nx.average_clustering(G):.2f}")
```

For this graph, G, we have L=1.800, C=0.60



In [7]: N = G.neighbors(0)
S = nx.subgraph(G, list(N))
nx.draw\_circular(S, \*\*opts)



- An (n,d)-circle graph has n nodes and m=nd edges.
- Each node has degree  $rac{2m}{n}=2d.$
- The social graph of each node has  $\displaystyle rac{3}{2} d(d-1)$  edges.
- The graph clustering coefficient of an (n, d)-circle graph is **independent of** n, and can be determined as

$$C=rac{3d-3}{4d-2}
ightarrowrac{3}{4} ext{, as }d
ightarrow\infty ext{.}$$

In particular:

### Characteristic Path Length

• However, things don't work as well when it comes to shortest paths (if we let  $n o \infty$ ). Indeed, the characteristic path length of an (n,d)-circle graph is approximately

$$L pprox rac{n}{4d},$$

growing linearly with n (for fixed d).

In conclusion, such regular graphs have **high clustering** but **long shortest paths**, hence (n, d)-circle graphs do not exhibit the small world behaviour.

To see how we could reduce the CPL, let's return to the Cycle Graph from earlier





So we can reduce the CPL, by adding relatively few edges. Finally, we can get a combined solution...

## The Watts-Strogatz Model

The following modification of the circle graph was suggested by Duncan J. Watts and Steven Strogatz (1998). The idea is to introduce a probabilistic element to the graph, which results in "shortcuts" (or "teleports") between the nodes and in a shortening of the characteristic path length.

**Definition (The WS Model).** Let 1 < d < n/2 and  $0 \le p \le 1$ . An (n, d, p)-WS graph G = (X, E) is constructed from an (n, d)-circle graph  $G_0 = (X, E_0)$  by rewiring each of the edges in  $E_0$  with probability p, as follows:

- 1. visit the nodes  $X=\{0,\ldots,n{-}1\}$  in turn ('clockwise').
- 2. for each node  $i \in X$  consider the d edges connecting i to j in a clockwise sense (  $j=i+1,\ldots,i+d$ ).
- 3. With probability p, in the edge (i,j) replace j by node  $k \in X$  chosen uniformly at random, subject to
  - k 
    eq i, and
  - (i,k) must not be an edge of G already.

```
In [12]: import random as rd
def ws_graph(n, d, p):
    G = circle_graph(n, d)
```

```
for v in G:
    for o in range(1, d+1):
        if rd.random() < p:
            w = rd.randint(0,n-1)  # pick a random node
            if w != v and not G.has_edge(v, w):
                G.remove_edge(v, (v+o) % n)
                G.add_edge(v, w)</pre>
```

**return** G

#### In [13]: n, d = 16, 3

```
G = ws_graph(n, d, 0.2)
nx.draw_circular(G, **opts)
print(f"G has L={nx.average_shortest_path_length(G):.3}, and C={nx.average_clustering
```

G has L=1.67, and C=0.42



In [14]: n, d = 16, 3
G = ws\_graph(n, d, 0.3)
nx.draw\_circular(G, \*\*opts)
print(f"G has L={nx.average\_shortest\_path\_length(G):.3}, and C={nx.average\_clustering
G has L=1.66, and C=0.49



In [15]: G = ws\_graph(n, d, 1)
nx.draw\_circular(G, \*\*opts)
print(f"G has L={nx.average\_shortest\_path\_length(G):.3} and C={nx.average\_clustering(C)}

G has L=1.66 and C=0.54



A WS graph with parameters (n, d, p) can be generated with the command:

```
nx.watts_strogatz_graph(n, 2*d, p).
```

```
In [16]: n, d = 21, 3
G = nx.watts_strogatz_graph(n, 2*d, 0.5)
nx.draw_circular(G, **opts)
print(f"G has L={nx.average_shortest_path_length(G):.3} and C={nx.average_clustering(C)}
```

G has L=1.75 and C=0.31  $\,$ 



In [17]: G = nx.watts\_strogatz\_graph(n, 2\*d, 0.1)
nx.draw\_circular(G, \*\*opts)
print(f"G has L={nx.average\_shortest\_path\_length(G):.3}, and C={nx.average\_clustering

G has L=1.94, and C=0.48  $\,$ 



In [18]: G = nx.watts\_strogatz\_graph(n, 2\*d, 0.2)
nx.draw\_circular(G, \*\*opts)
print(f"G has L={nx.average\_shortest\_path\_length(G):.3}, C={nx.average\_clustering(G):

G has L=1.89, C=0.45



## **Properties of WS-Graphs**

The small-world attributes of a (n, d, p)-WS graph depend on the probability p. The following measurements have been taken for n = 1000 and d = 5.

p	L	C
0	50.5	0.667
0.01	8.94	0.648
0.05	5.26	0.576
1	3.27	0.00910

## Exercises

- 1. In terms of the parameters, n, d and p, what is the clustering coefficient C of an (n, d, p)-WS graph?
- 2. In terms of the parameters, n, d and p, what is the average shortest path length L of an (n, d, p)-WS graph?