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## CS4423-Networks: Week 10 (19+20 March 2025)

### Part 2: Clustering

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This Jupyter notebook, and PDF and HTML versions, can be found at <https://www.niallmadden.ie/2425-CS4423/#Week10>

*This notebook was written by Niall Madden, adapted from notebooks by Angela Carnevale.*

#### Modules for this notebook

```
In [1]: import networkx as nx
import numpy as np
opts = { "with_labels": True, "node_color": "#004225", "font_color": "white" } # Iris

import random # some random number generators: random, random_choices
import statistics # e.g., mean of entries in a list
import math # for comb (=binomial coef)
import matplotlib.pyplot as plt

np.set_printoptions(precision=2) # just display arrays to 2 decimal places
np.set_printoptions(suppress=True)
```

### Giant Components (again)

In Part 1, we learned about the following result:

Suppose  $p(n) = cn^{-1}$  for some positive constant  $c$ . (Then the average degree  $\langle k \rangle = pn = c$  remains fixed as  $n \rightarrow \infty$ .)

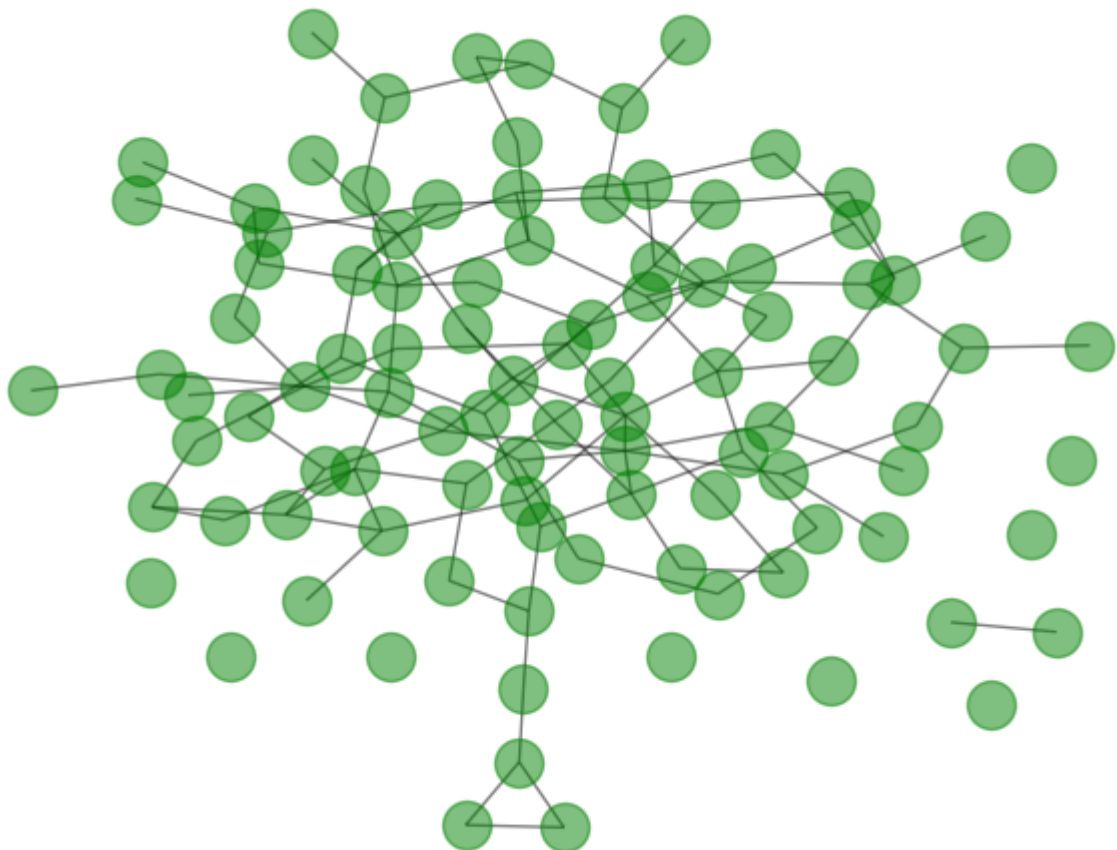
**Theorem (Erdős-Rényi).**

- If  $c < 1$  the graph contains many small components, orders bounded by  $O(\ln n)$ .
- If  $c = 1$  the graph has large components of order  $O(n^{2/3})$ .
- If  $c > 1$  there is a unique **giant component** of order  $O(n)$ .

There are a few ways we can verify this. One is to fix  $n$ , and construct a graph for values of  $p$  that correspond to these three ranges:

```
In [2]: n=100
p=2/n # take p=0.5/n, 1/n, and 2/n
G = nx.gnp_random_graph(n, p, seed=6)
pos = nx.nx_agraph.graphviz_layout(G, prog="neato", args="")
nx.draw(G, pos, alpha=0.5, node_color="green", with_labels=False)
S = len(max(nx.connected_components(G), key=len))
print(f"p={p}, largest component has {S} nodes")
```

p=0.02, largest component has 89 nodes



Another way, is to compute the order the largest component for a range of values of  $c$ . Note: this test might take about 30 seconds to run.

```
In [3]: n = 1000 # order of the graphs
c = np.linspace(0, 3, 31) # range of values of c
Ps = (1/n)*c # corresponding probabilities
```

```

S = []
Runs = 5      # Average over 5 graphs
for p in Ps:
    Ave = 0
    for i in range(Runs):
        G = nx.gnp_random_graph(n,p)
        Ave += len(max(nx.connected_components(G), key=len))
    S += [Ave/Runs]

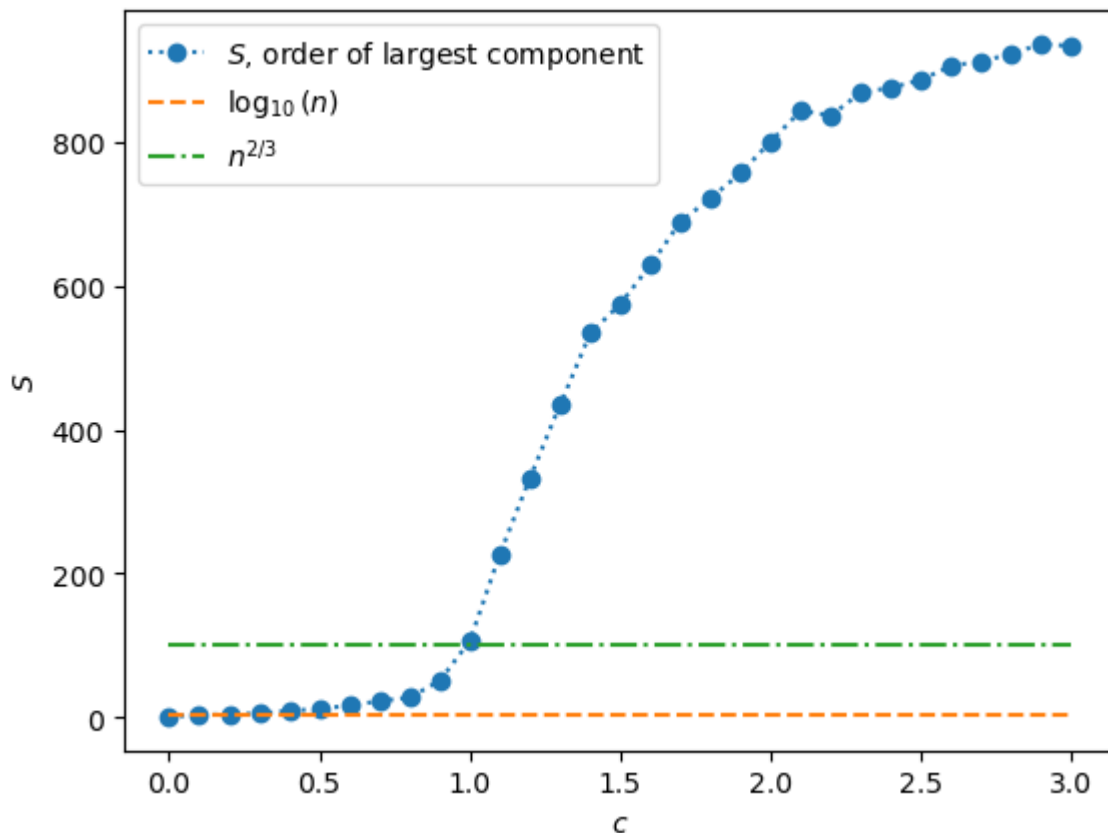
```

```

In [4]: plt.plot(c, S, ':o', label=r'$S$, order of largest component')
plt.plot(c, c*0 + math.log10(n), '--', label=r'$\log_{10}(n)$')
plt.plot(c, c*0 + n**(2/3), '-.', label=r'$n^{2/3}$')
plt.xlabel(r'$c$')
plt.ylabel(r'$S$')
plt.legend()

```

Out[4]: <matplotlib.legend.Legend at 0x7f92c81d2f90>



## Characteristic Path Length

We learned in Part 1 that there are several measures of "Small Worldedness", that include

- Small Characteristic Path Length (**CPL**)
- High Clustering (not yet defined)

We asserted in that class that graphs in the  $G_{ER}$  models tend to have small Characteristic Path Length, but not high clustering.

Recall: the **characteristic path length**  $L$  of  $G$  is the average distance between pairs of distinct nodes,

$$L = \frac{1}{n(n-1)} \sum_{i \neq j} d_{ij}.$$

## CLP in $G_{ER}(n, m)$

**Fact** (noted in Part 1): The characteristic path length of a random network in  $G_{ER}(n, m)$ , or  $G_{ER}(n, p)$  is

$$L = \frac{\ln n}{\ln \langle k \rangle}.$$

where  $\langle k \rangle$  is the average degree of the network.

Consider  $G_{ER}(n, m)$ , where we know  $\langle k \rangle = 2m/n$ .

So if  $n = 16$  and  $m = 32$ , then the average node degree in  $G(n, m)$  is  $\langle k \rangle = \frac{2m}{n} = 4$ , and, approximately,  $L = \frac{\log_2 16}{\log_2 4} = 2$ .

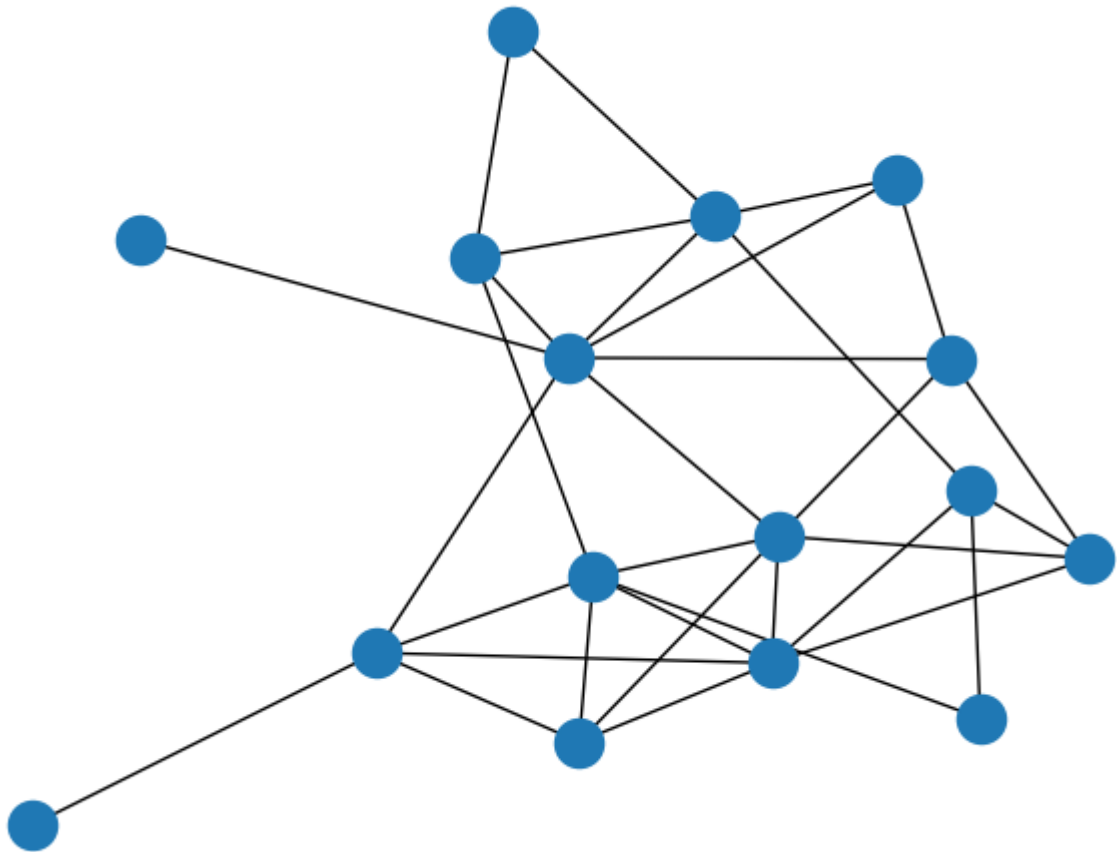
In `networkx`, this is called `average_shortest_path_length`. However, it can only be applied to connected graphs. So, let's choose  $G$  from  $G_{ER}(16, 32)$  a connected graph:

```
In [5]: n = 16
m = 32
G = nx.gnm_random_graph(n,m)
count=0
while (not nx.is_connected(G)):
    count +=1
    G = nx.gnm_random_graph(n,m)
print(f"FYI, took {count} iterations to get a connected G")
L = nx.average_shortest_path_length(G)
print(f"G_ER({n},{m}) has an average shortest path length of {L:.3f}")
```

```

FYI, took 1 iterations to get a connected G
G_ER(16,32) has an average shortest path length of 2.017
```

```
In [6]: nx.draw(G)
```



## Computing CPL by hand

Yesterday, we learned about CPL in the context of eccentricity, in turn computed from by the distance matrix  $\mathcal{D}$ . We can compute  $\mathcal{D}$  for  $G$  as:

```
In [7]: dist = dict(nx.shortest_path_length(G))
Dm = [[dist[i][j] for j in range(n)] for i in range(n)]
print(np.array(Dm))
```

```
[[0 3 3 3 2 3 3 3 2 1 3 3 4 2 3 2]
 [3 0 2 2 2 2 2 2 1 2 2 1 1 1 3 2]
 [3 2 0 2 1 2 2 2 1 2 1 2 3 1 1 1]
 [3 2 2 0 2 1 4 3 1 2 3 2 3 3 2 2]
 [2 2 1 2 0 1 3 3 1 1 2 2 3 2 1 1]
 [3 2 2 1 1 0 3 2 2 2 2 1 2 2 1 2]
 [3 2 2 4 3 3 0 2 3 2 2 2 3 1 3 3]
 [3 2 2 3 3 2 2 0 3 2 1 1 2 1 2 3]
 [2 1 1 1 1 2 3 3 0 1 2 2 2 2 2 1]
 [1 2 2 2 1 2 2 2 1 0 2 2 3 1 2 1]
 [3 2 1 3 2 2 2 1 2 2 0 2 3 1 1 2]
 [3 1 2 2 2 1 2 1 2 2 2 0 1 1 2 3]
 [4 1 3 3 3 2 3 2 2 3 3 1 0 2 3 3]
 [2 1 1 3 2 2 1 1 2 1 1 1 2 0 2 2]
 [3 3 1 2 1 1 3 2 2 2 1 2 3 2 0 2]
 [2 2 1 2 1 2 3 3 1 1 2 3 3 2 2 0]]
```

As we know, the **characteristic path length**  $L$  is the sum of all entries in  $\mathcal{D}$ , divided by the number of pairs of distinct nodes  $n(n-1)$ .

```
In [8]: cpl = sum([sum(d) for d in Dm])/(n*(n-1))
```

```
print(f"CPL = {cpl:.3f}")
```

CPL = 2.017

## Small World

**Definition (Small-world behaviour).** A network  $G = (X, E)$  is said to exhibit a **small world behaviour** if its characteristic path length  $L$  grows proportionally to the logarithm of the number  $n$  of nodes of  $G$ :

$$L \sim \ln n.$$

In this sense, the ensembles  $G(n, m)$  and  $G(n, p)$  of random graphs do exhibit small world behavior (as  $n \rightarrow \infty$ ).

## Transitivity

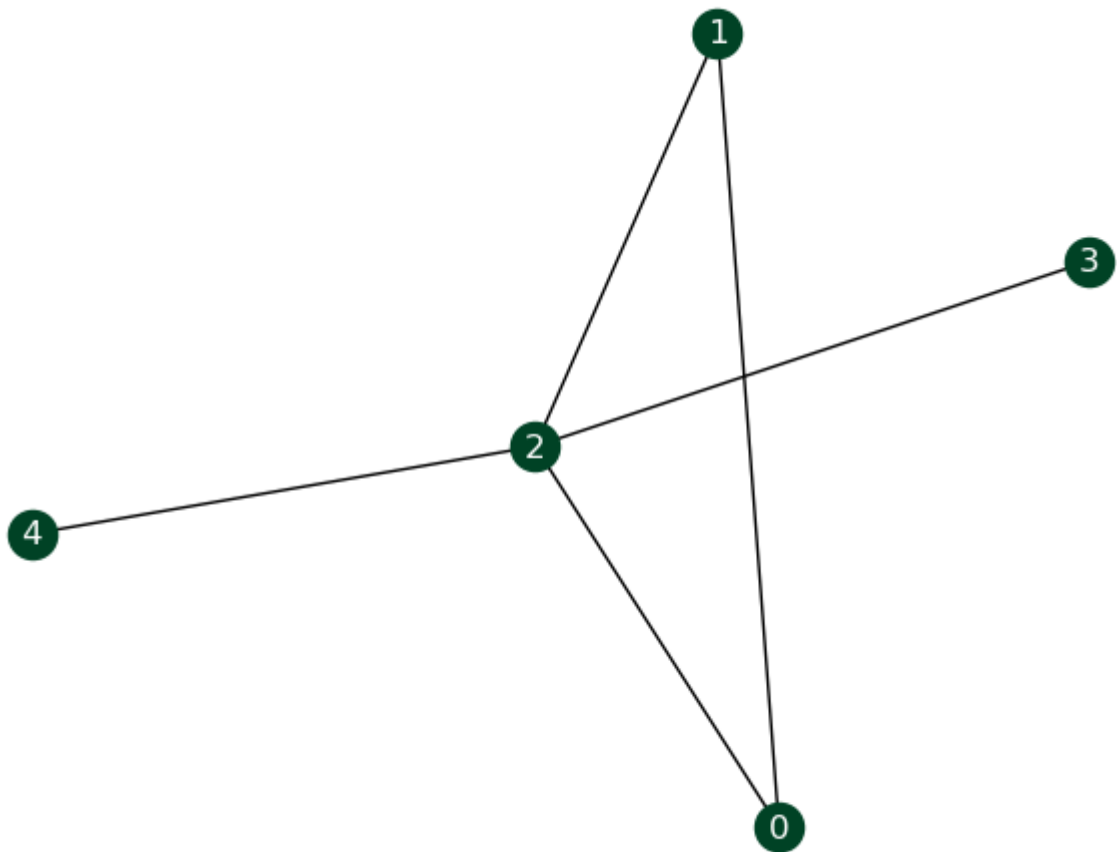
**Definition (Graph transitivity).** The **transitivity**  $T$  of a graph  $G = (X, E)$  is the proportion of **transitive** triads, i.e., triads which are subgraphs of **triangles**. This proportion can be computed as follows:

$$T = \frac{3n_{\Delta}}{n_{\wedge}},$$

where  $n_{\Delta}$  is the number of triangles in  $G$ , and  $n_{\wedge}$  is the number of triads.

Example:

```
In [9]: G = nx.Graph(((0,1), (1,2), (2,0), (2,3), (2,4)))
        nx.draw(G, **opts)
```



The function `nx.triangles(G)` returns a python dictionary reporting for each node of the graph `G` the number of triangles it is contained in.

```
In [10]: print(nx.triangles(G))
```

```
{0: 1, 1: 1, 2: 1, 3: 0, 4: 0}
```

Overall, each triangle in `G` is thus accounted for 3 times, once for each of its nodes. Hence, the following sum determines this number  $3n_{\Delta}$ .

```
In [11]: NumTriangles = sum(nx.triangles(G).values())/3
print(f"G has {NumTriangles} triangle(s)")
```

```
G has 1.0 triangle(s)
```

As we've seen, the number  $n_{\Delta}$  of triads in `G` can be determined from the graph's degree sequence, as each node of degree  $k$  is the central node of exactly  $\binom{k}{2}$  triads.

```
In [12]: NumTriads = 0
for i in G.nodes():
    NumTriads += math.comb(G.degree[i],2)
print(f"G has {NumTriads} triade(s)")
```

```
G has 8 triade(s)
```

```
In [13]: T = 3*NumTriangles/NumTriads
print(f"Transitivity of G is {T}")
```

```
Transitivity of G is 0.375
```

Of course, there is a built-in function to do this:

```
In [14]: nx.transitivity(G)
```

```
Out[14]: 0.375
```

## Transitivity of $G_{ER}(n, p)$

The transitivity of graph in  $G_{ER}(n, p)$  is easy to estimate: for every triad, the "third" edge is present with probability  $p$ . So

$$T = p,$$

(Or: Compute  $3n_{\Delta}/n_{\wedge}$  using the explicit formulas from the previous lecture:  $n_{\Delta} = \binom{n}{3}p^3$  and  $n_{\wedge} = 3\binom{n}{3}p^2$ .)

Let's check:

```
In [15]: n,p = 100, 0.1
T=nx.transitivity(nx.gnp_random_graph(n,p))
print(f"G({n},{p}) has T = {T:.3}")
```

```
G(100,0.1) has T = 0.103
```

## Clustering

The concept of **clustering** measures the transitivity of a node, or of an entire graph in a different way.

To define it, we need the concept of an **induced subgraph**.

### Induced subgraph

Given  $G = (X, E)$  and  $Y \subset X$ , the induced subgraph of  $G$  on  $Y$  is the graph  $H = \left(Y, E \cap \binom{Y}{2}\right)$ .

That is:

- $H$  is a subgraph of  $G$ , with node set  $Y$
- $H$  has all possible edges in  $G$  for which both nodes are in  $Y$ .

(See examples on board).

In `networkx`, we can get an induced subgraph of  $G$ , on nodes  $\{x, y, z\}$ , from  
`G.subgraph([x,y,z])`

## Clustering coefficient

**Definition (Clustering coefficient).** For a node  $i \in X$  of a graph  $G = (X, E)$ , denote by  $G_i$  the subgraph induced on the neighbours of  $i$  in  $G$ , and by  $m(G_i)$  its number of edges.

The **node clustering coefficient**  $c_i$  of node  $i$  is defined as

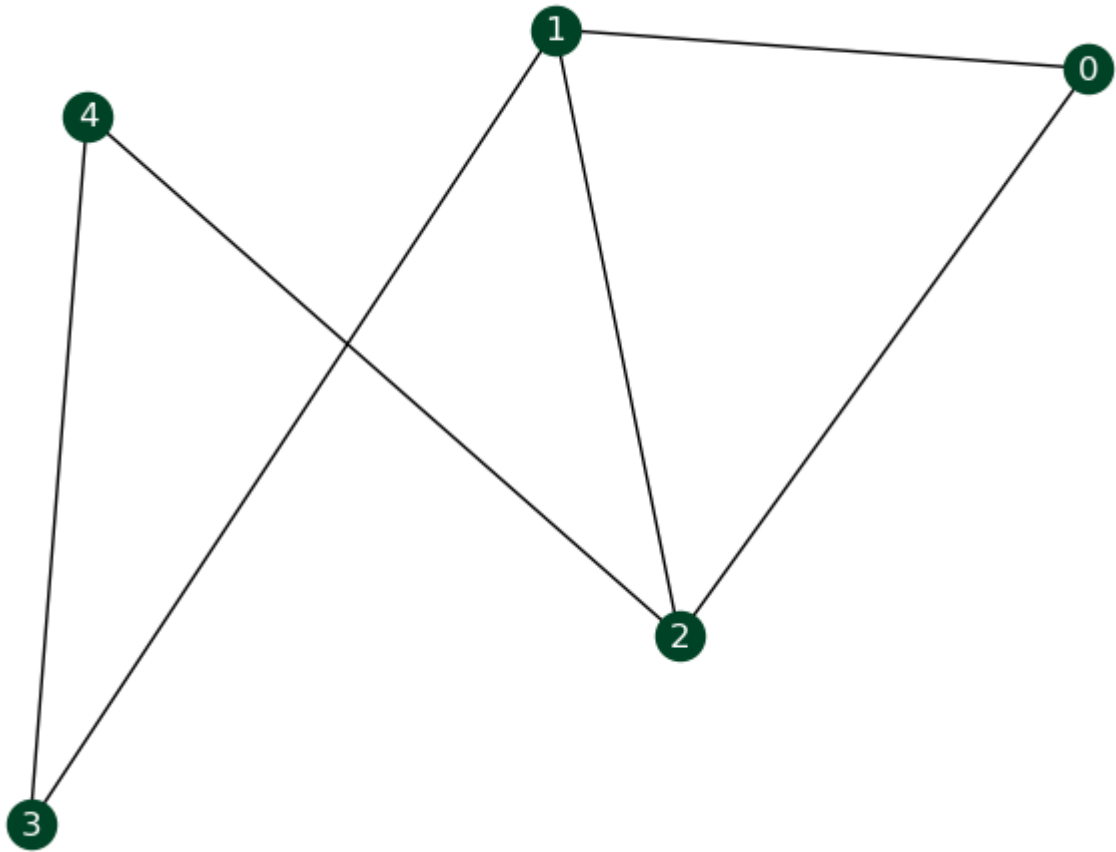
$$c_i = \begin{cases} \binom{k_i}{2}^{-1} m(G_i), & k_i \geq 2, \\ 0, & \text{else.} \end{cases}$$



That is, the node clustering coefficient measures the proportion of existing edges its **social graph** among the possible edges.

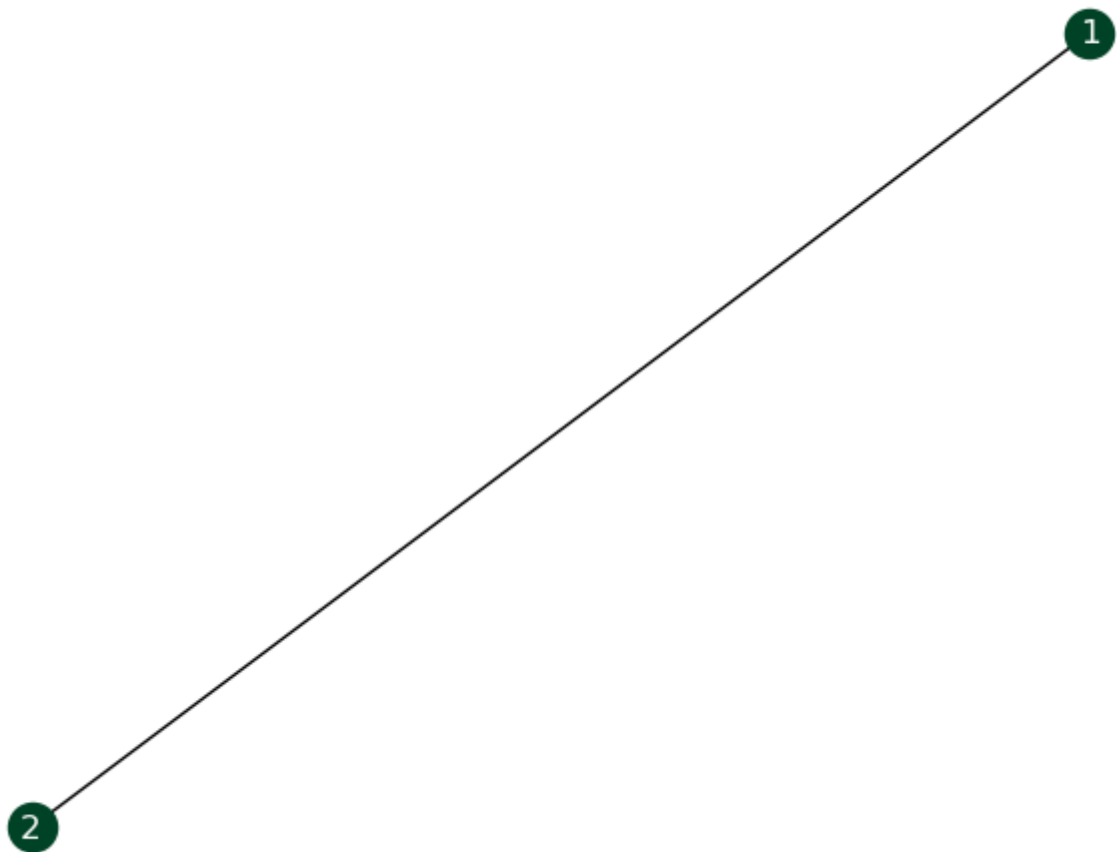
**Example:**

```
In [16]: G = nx.Graph([(0,1), (0,2), (1,2), (1,3), (2,4), (3,4)])  
         nx.draw(G, **opts)
```



The graph induced by the neighbours of 0 is :

```
In [17]: N = nx.neighbors(G, 0)  
         S = G.subgraph(list(N))  
         nx.draw(S, **opts)
```



Calculate the clustering coefficient for the nodes of  $G$ :

```
In [18]: for i in G.nodes():
         k_i = G.degree(i)
         N = nx.neighbors(G, i)
         S = G.subgraph(list(N))
         c_i = S.size()/math.comb(k_i,2)
         print(f"Node {i} has clustering coef {c_i:0.4}")
```

```
Node 0 has clustering coef 1.0
Node 1 has clustering coef 0.3333
Node 2 has clustering coef 0.3333
Node 3 has clustering coef 0.0
Node 4 has clustering coef 0.0
```

Of course, there is a `networkx` function for this:

```
In [19]: nx.clustering(G)
```

```
Out[19]: {0: 1.0, 1: 0.3333333333333333, 2: 0.3333333333333333, 3: 0, 4: 0}
```

## Graph Clustering Coefficient

The **graph clustering coefficient**  $C$  of  $G$  is the average node clustering coefficient,

$$C = \langle c \rangle = \frac{1}{n} \sum_{i=1}^n c_i.$$

By definition,  $0 \leq c_i \leq 1$  for all nodes  $i \in X$ , and  $0 \leq C \leq 1$ .

```
In [20]: nx.average_clustering(G)
```

```
Out[20]: 0.3333333333333333
```

## Clustering for $G_{ER}(n, p)$

The **node clustering coefficient** of any node  $i$  in a  $G_{ER}(n, p)$  **random graph** is  $c_i = p$ . (In any selection of potential edges, by construction a proportion  $p$  of them is present in the random graph; this is true in particular for the  $\binom{k}{2}$  potential edges between the  $k$  neighbors of a node of degree  $k$ .)

Thus the **graph clustering coefficient** of a  $G_{ER}(n, p)$  **random graph** is

$$C = p.$$

## IMPORTANT: Large $G_{ER}$ graphs have few triangles

Note that when  $p(n) = \langle k \rangle n^{-1}$  for a fixed expected average degree  $\langle k \rangle$  then  $C = \langle k \rangle / n \rightarrow 0$  for  $n \rightarrow \infty$ . That is **in large  $G_{ER}$  random graphs, the number of triangles is negligible**.

In real world networks, one often observes that  $C / \langle k \rangle$  does not depend on  $n$  (as  $n \rightarrow \infty$ ).

## Clustering vs Transitivity

For a node  $i \in X$ , denote by  $n_i^\wedge = \binom{k_i}{2}$  the number of triads containing  $i$  as their central node, and by  $n_i^\Delta$  the actual number of triangles containing  $i$ .

Then the node clustering coefficient is  $c_i = n_i^\Delta / n_i^\wedge$ , or  $n_i^\Delta = n_i^\wedge c_i$ .

Moreover  $3n_\Delta = \sum_i n_i^\Delta$  and  $n_\wedge = \sum_i n_i^\wedge$ .

It follows that

$$T = \frac{3n_\Delta}{n_\wedge} = \frac{1}{n_\wedge} \sum_i n_i^\wedge c_i$$

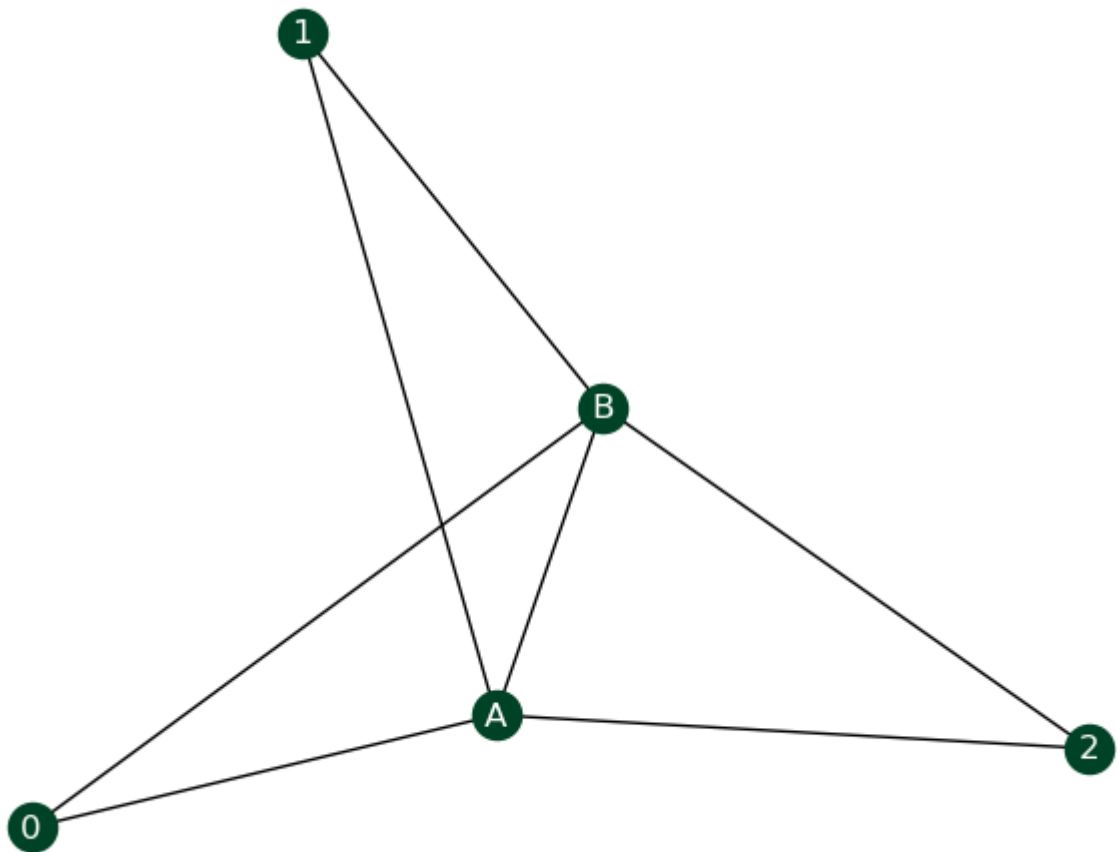
in contrast to

$$C = \frac{1}{n} \sum_i c_i.$$

That is,  $C$  is the (plain) **average** of the node clustering coefficients, whereas  $T$  is a **weighted average** of node clustering coefficients, giving higher weight to high degree nodes.

The following example illustrates how  $C$  and  $T$  are different measures. This (very non-random) ensemble of networks has the property that, as  $n \rightarrow \infty$  here,  $T \rightarrow 0$  while  $C \rightarrow 1$ .

```
In [21]: n = 3 # try larger and larger n
G = nx.Graph(["AB"])
G.add_edges_from([(x, k) for x in "AB" for k in range(n)])
nx.draw(G, **opts)
```



```
In [22]: nx.average_clustering(G), nx.transitivity(G)
```

```
Out[22]: (0.8, 0.6)
```

- The fact that ER random networks tend to have low transitivity and clustering shows the need of a new kind of (random) network construction that is better at modelling real world networks.
- One idea, developed by Watts and Strogatz in 1998, is to start with some **regular network** that naturally has a **high clustering**, and then to randomly distort its edges, to introduce some **short paths**.

**FINISHED HERE THURSDAY**